



Lightweight Session Programming in Scala

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Troubles with session programming

Consider a simple “greeting” client/server session protocol:

1. the client can ask to **greet** someone, or **quit**
2. *if asked to greet*, the server can either:
 - 2.1 say **hello**, and go **back to 1**
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Typical approach:

- ▶ describe the protocol **informally**
- ▶ develop *ad hoc* **protocol APIs** to avoid **protocol violations**
- ▶ find bugs via **runtime testing/monitoring**

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Impact on **software evolution and maintenance**

Lightweight Session Programming in Scala

This talk: we show how in **Scala** + `lchannels` we can write:

```
def client(c: Out[Start]): Unit = {
  if (Random.nextBoolean()) {
    val c2 = c !! Greet("Alice")_

    c2 ? {
      case m @ Hello(name) => client(m.cont)
      case Bye(name)       => ()
    }
  } else {
    c ! Quit()
  }
}
```

...with a **clear theoretical basis**, giving a **general API** with **static protocol checks** and **message transport abstraction**



- ▶ **Object-oriented *and* functional**
- ▶ **Declaration-site variance**
- ▶ **Case classes** for OO pattern matching



- ▶ **Object-oriented** *and* **functional**
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```
sealed abstract class Pet
case class Cat(name: String) extends Pet
case class Dog(name: String) extends Pet
```

```
def says(pet: Pet) = {
  pet match {
    case Cat(name) => name + " says: meow"
    case Dog(name) => name + " says: woof"
  }
}
```

Session types

Consider again our **“greeting” client/server session protocol**:

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We can **formalise** the **client** viewpoint as a **session type** for the **session π -calculus**: (Honda *et al.*, 1993, 1994, 1998, ...)

$$S_h = \mu X. \left(\begin{array}{l} !\text{Greet}(\text{String}). \left(\begin{array}{l} ?\text{Hello}(\text{String}). X \\ \& \\ ?\text{Bye}(\text{String}). \text{end} \end{array} \right) \\ \oplus \\ !\text{Quit}. \text{end} \end{array} \right)$$

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We can **formalise** the **server** viewpoint as a (*dual*) **session type** for the **session** π -calculus: (Honda *et al.*, 1993, 1994, 1998, ...)

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From theory to practice

Desiderata:

- ▶ find a **formal link** between **Scala types** and **session types**
- ▶ represent **sessions** in a language **without session primitives**
 - ▶ **lightweight**: no language extensions, minimal dependencies

Inspiration (from concurrency theory):

- ▶ **encoding of session types into linear types for π -calculus**
(Dardha, Giachino & Sangiorgi, PPDP'12)

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Result: **Lightweight Session Programming in Scala**

Session vs. linear types (in pseudo-Scala)

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```

“Linear Scala”

```
def client(c: LinOutChannel[?]): Unit = {
  if (...) {
    val (c2in, c2out) = createLinChannels[?]()
    c.send( Greet("Alice", c2out) )
    c2in.receive match {
      case Hello(name, c3out) => client(c3out)
      case Bye(name)         => ()
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Goals:

- ▶ define and implement linear in/out channels
- ▶ instantiate the “?” type parameter
- ▶ automate continuation channel creation

Ichannels: interface

```
abstract class In[+A] {  
  
  def receive(implicit d: Duration): A  
  
}  
  
abstract class Out[-A] {  
  
  def send(msg: A): Unit  
  
}
```

API offers **typed** `send/receive`

- ▶ with **runtime checks** for **linear use** and **error handling**

Note **input/output co/contra-variance**

lchannels: interface

```

abstract class In[+A] {

  def receive(implicit d: Duration): A

  def ?[B](f: A => B)(implicit d: Duration): B = {
    f(receive)
  }
}

abstract class Out[-A] {

  def send(msg: A): Unit
  def !(msg: A)                = send(msg)
}

```

API offers **typed** `send/receive`, plus **syntactic sugar**

- ▶ with **runtime checks** for **linear use** and **error handling**

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Ichannels: interface

```

abstract class In[+A] {
  def future: Future[A]
  def receive(implicit d: Duration): A = {
    Await.result[A](future, d)
  }
  def ?[B](f: A => B)(implicit d: Duration): B = {
    f(receive)
  }
}

abstract class Out[-A] {
  def promise[B <: A]: Promise[B] // Impl. must be constant
  def send(msg: A): Unit          = promise.success(msg)
  def !(msg: A)                   = send(msg)
}

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    f(receive)
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abstract class Out[-A] {
  def promise[B <: A]: Promise[B] // Impl. must be constant
  def send(msg: A): Unit          = promise.success(msg)
  def !(msg: A)                   = send(msg)
  def create[B]() : (In[B], Out[B]) // Used to continue a session
}

```

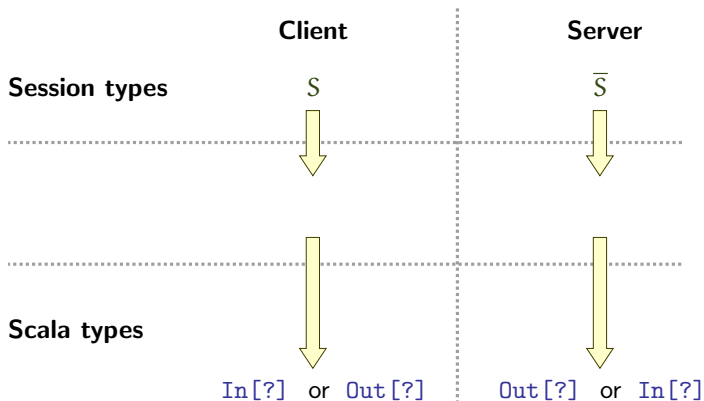
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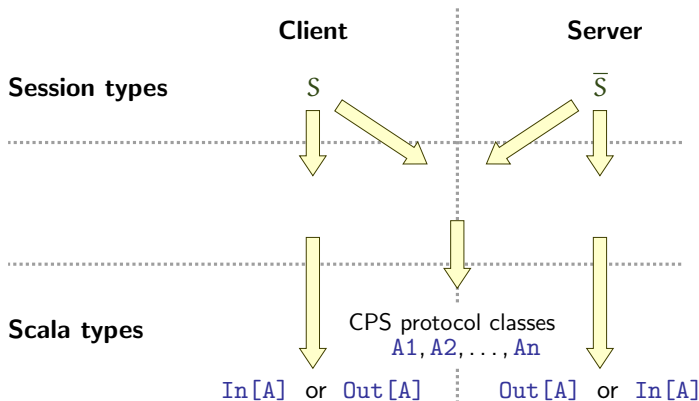
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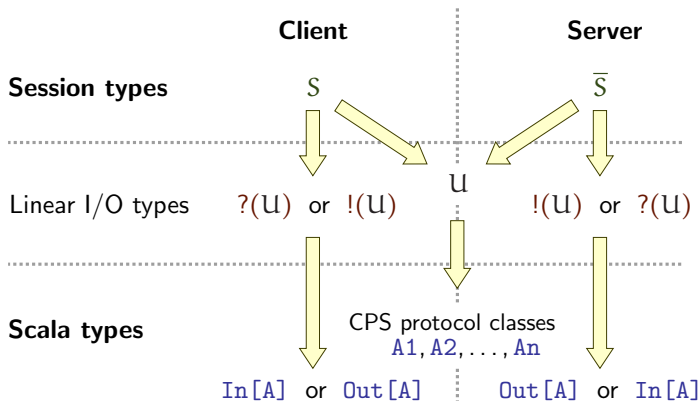
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Programming with lchannels (I)

$$S_h = \mu_X. \left(!\text{Greet}(\text{String}). (? \text{Hello}(\text{String}). X \ \& \ ? \text{Bye}(\text{String}). \text{end}) \oplus !\text{Quit}. \text{end} \right)$$

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```
// Top-level internal choice
case class Greet(p: String)
case class Quit(p: Unit)
```

$\text{prot} \ll S_h \gg_{\mathcal{N}} =$

```
// Inner external choice
case class Hello(p: String)
case class Bye(p: String)
```

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sealed abstract class Start
case class Greet(p: String)           extends Start
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case class Hello(p: String)          extends Greeting
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case class Greet(p: String)(val cont: Out[Greeting]) extends Start
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Goals:

- define and implement linear in/out channels ✓
- instantiate the "?" type parameter ✓
- automate continuation channel creation ✗

The “create-send-continue” pattern

We can observe that `In/Out` channel pairs are usually created for **continuing a session after sending a message**

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Let us **add the `!!` method** to `Out[·]`:

```
abstract class Out[-A] {  
  ...  
  def !![B](h: Out[B] => A): In[B] = {  
    val (cin, cout) = this.create[A]() // Create...  
    this ! h(cout) // ...send...  
    cin // ...continue  
  }  
  
  def !![B](h: In[B] => A): Out[B] = {  
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Scala + lchannels

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Demo

Run-time and compile-time checks

Well-typed output / int. choice
Exhaustive input / ext. choice

Compile-time
Compile-time

Run-time and compile-time checks

Well-typed output / int. choice
Exhaustive input / ext. choice

Compile-time
Compile-time

Double use of linear output endp.
Double use of linear input endp.

Run-time
Run-time

Run-time and compile-time checks

Well-typed output / int. choice

Exhaustive input / ext. choice

Compile-time

Compile-time

Double use of linear output endp.

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Run-time

Run-time

“Forgotten” output

“Forgotten” input

Run-time (timeout on input side)

Unchecked

Formal properties

Theorem (*Preservation of duality*).

$$\langle\langle \overline{S} \rangle\rangle_{\mathcal{N}} = \overline{\langle\langle S \rangle\rangle_{\mathcal{N}}} \quad (\text{where } \overline{\text{In}[A]} = \text{Out}[A] \text{ and } \overline{\text{Out}[A]} = \text{In}[A]).$$

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Theorem (*Dual session types have the same CPS protocol classes*).

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$$\text{prot}\langle\langle S \rangle\rangle_{\mathcal{N}} = \text{prot}\langle\langle \bar{S} \rangle\rangle_{\mathcal{N}}.$$

Theorem (*Scala subtyping implies session subtyping*).

For all S, \mathcal{N} :

- ▶ if $\langle\langle S \rangle\rangle_{\mathcal{N}} = \text{In}[A]$ and $B <: \text{In}[A]$,
then $\exists S', \mathcal{N}'$ such that $B = \langle\langle S' \rangle\rangle_{\mathcal{N}'}$, and $S' \leq S$;
- ▶ if $\langle\langle S \rangle\rangle_{\mathcal{N}} = \text{Out}[A]$ and $\text{Out}[A] <: B$,
then $\exists S', \mathcal{N}'$ such that $B = \langle\langle S' \rangle\rangle_{\mathcal{N}'}$, and $S \leq S'$.

Conclusions

We presented a **lightweight integration of session types in Scala** based on a **formal link** between CPS protocols and session types

We leveraged **standard Scala features** (from its type system and library) with a **thin abstraction layer** (`lchannels`)

- ▶ low **cognitive overhead, integration** and **maintenance** costs
- ▶ naturally supported by **modern IDEs** (e.g. **Eclipse**)

We validated our session-types-based programming approach with **case studies** (from literature and industry) and **benchmarks**

Ongoing and future work

Automatic generation of CPS protocol classes
from session types, using **Scala macros**

- ▶ *B. Joseph. "Session Metaprogramming in Scala". MSc Thesis, 2016*

Extension to **multiparty session types**, using **Scribble**

- ▶ A. Scalas, O. Dardha, R. Hu, N. Yoshida.
*"A Linear Decomposition of Multiparty Sessions
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Generalise the approach to other frameworks beyond

[lchannels](#), and study its properties.

Natural candidates: **Akka Typed**, **Reactors.IO**

Investigate other programming languages. Possible candidate:

C# (declaration-site variance and FP features)

Try lchannels and Scribble!

<http://alcestes.github.io/lchannels>

<http://scribble.org>



ECOOP 2016



ECOOP 2017